

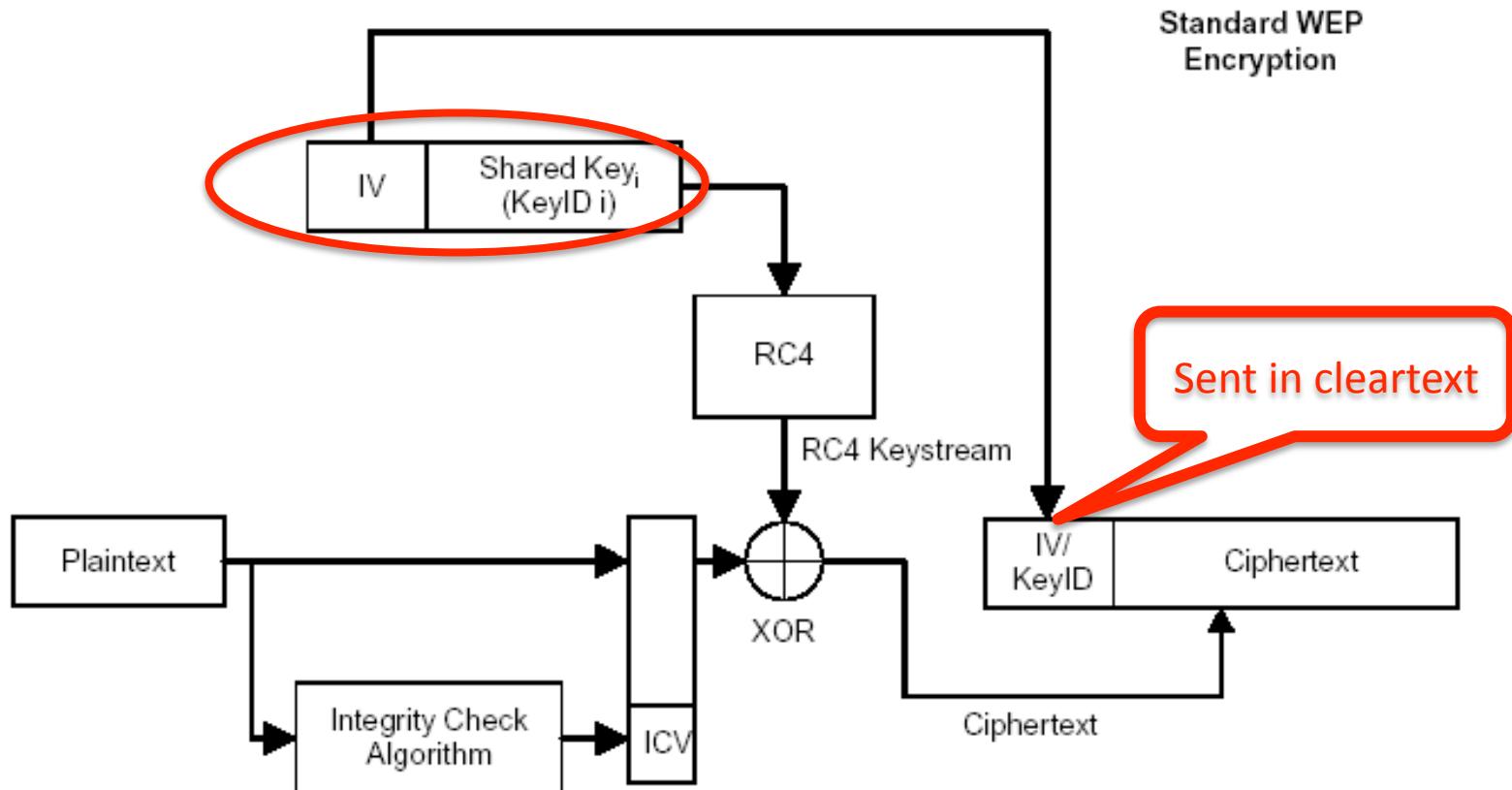


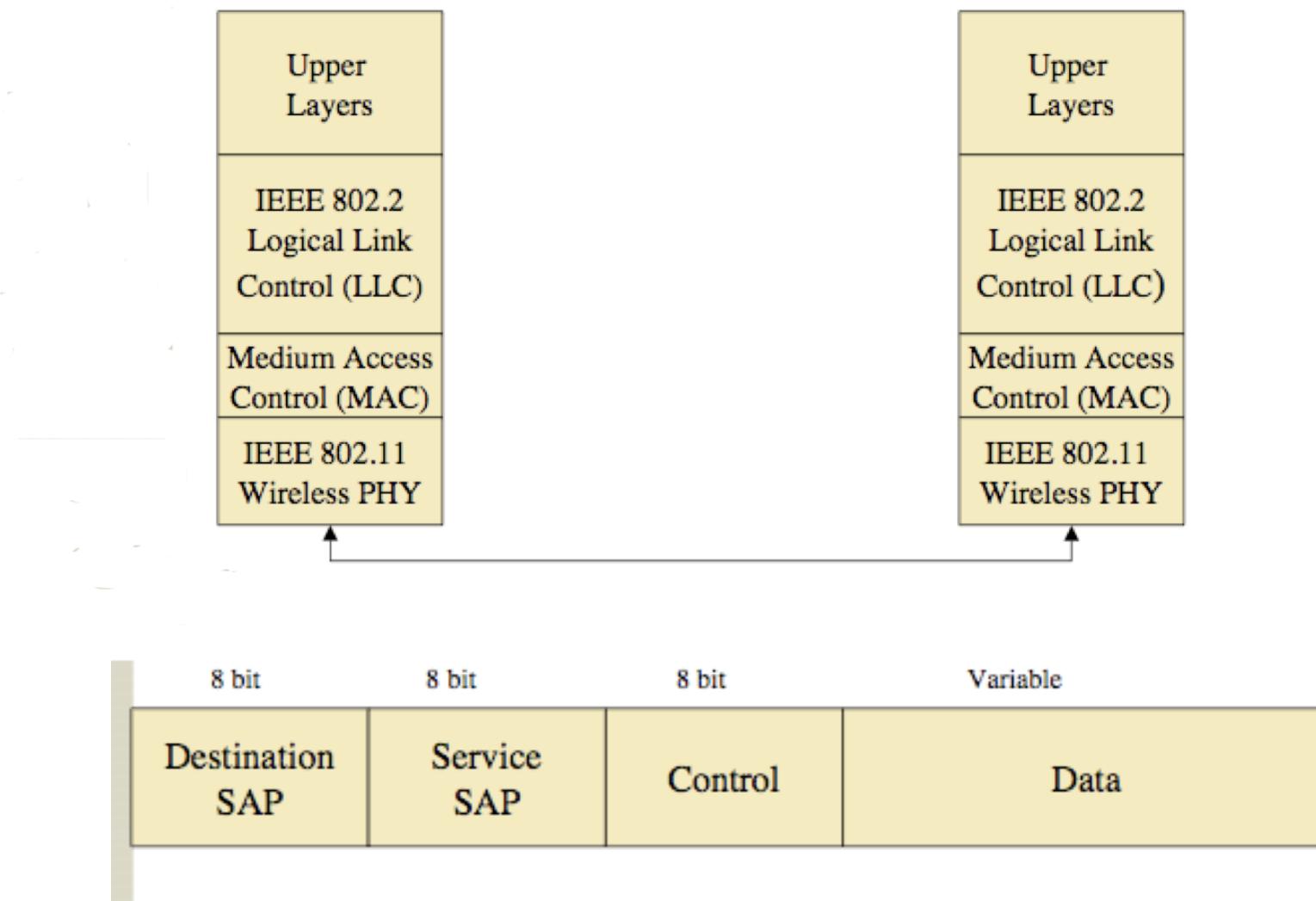
# WEP

- Goals
  - Authentication
    - A pre-shared key K, shared among all APs and STAs members of the WLAN
    - Same key used for authentication and confidentiality
  - Confidentiality
    - RC4
  - Integrity Protection
    - CRC-32 (non-crypto primitive)

# RC4

Figure 1. Standard WEP Encryption Process





- RC4 uses the key and the IV to initialize a state machine via `ksa()`
- RC4 modifies the state and generates a new byte of the keystream from the new state using `prga()`;
- first byte of the plaintext comes from the WEP SNAP header,
  - the first byte of the keystream can be obtained from

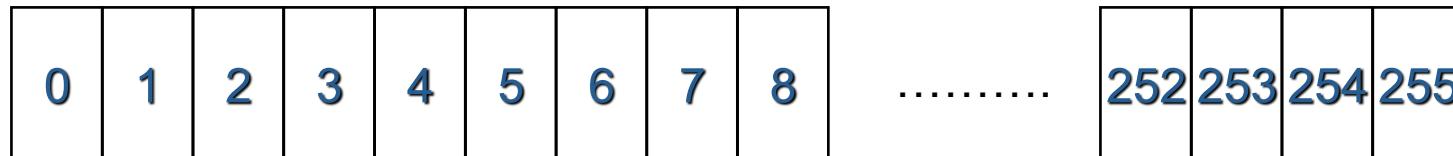
$$B \oplus 0xAA$$

- We can simulate the start of KSA because IV is sent in cleartext and  $K = IV | K$
- Attacker knowing a byte of the keystream can derive a byte of the key (weakness in PRNG)
- Weak IV:  $(A+3, n-1, X)$ 
  - Initialization vector that with 5% of probability allow to recover a byte of the key
  - The byte of keystream is linked to the value of the key

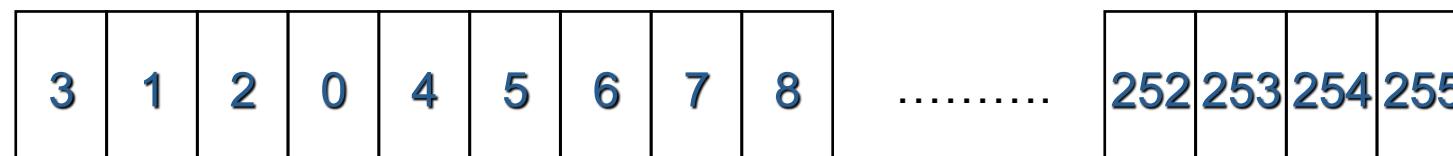
# KSA

- key scheduling algorithm (KSA).
  - begin ksa(with int keylength, with byte key[keylength])
    - for i from 0 to 255
      - $S[i] := i$
    - Endfor
    - $j := 0$
    - for i from 0 to 255
      - $j := (j + S[i] + \text{key}[i \bmod \text{keylength}]) \bmod 256$
      - swap( $S[i], S[j]$ )
    - Endfor
    - end

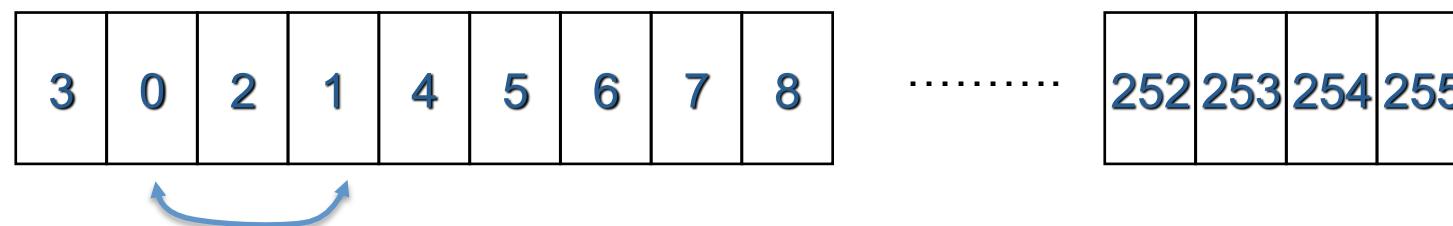
# KSA: scrambling the byte array



Step 1



Step 2



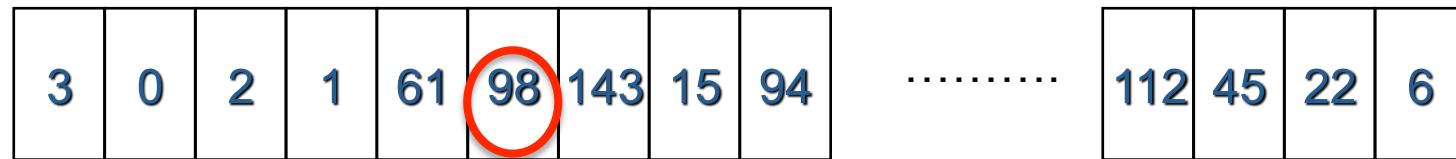
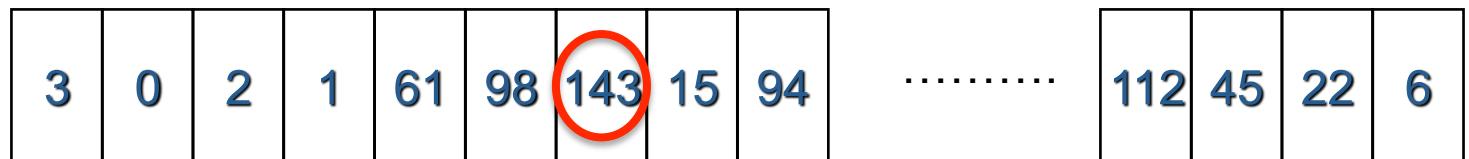
Step 3

$\text{IV} = 3,255,X$

# PRGA

- pseudo-random generation algorithm (PRGA)
  - begin prga(with byte S[256])
    - i := 0
    - j := 0
    - while GeneratingOutput:
      - i := (i + 1) mod 256
      - j := (j + S[i]) mod 256
      - swap(S[i],S[j])
      - output S[(S[i] + S[j]) mod 256]
    - Endwhile
  - end

# PRGA: extract a byte of keystream



- If  $S$  is a permutation,  $S^{-1}$  is the inverse permutation
- If an attacker knows the first  $l$  bytes of an RC4 key used to generate a keystream  $X()$ .
  - He can simulate the first  $l$  steps of the RC4-KSA ( $S_l$  and  $j_l$  known)
  - $S_l[1] < l$
  - $S_l[1] + S_l[S_l[1]] = l$
  - $S^{-1}_l[X[0]] \neq 1$
  - $S^{-1}_l[X[0]] \neq S_l[1]$
- The function  $F()$  will take the value of  $K[l]$  with a probability of about 5%
  - $F(K[0], \dots, K[l-1], X[0]) = S^{-1}[X[0]] - j - S[l]$

# WPA-PSK

**[Message 1: A → S]**

AA, ANonce, sn, msg1

**[Message 2: S → A]**

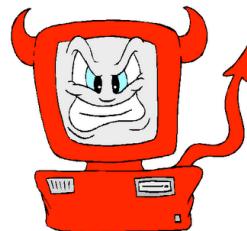
SPA, SNonce, sn, msg2,  $\text{MIC}_{\text{PTK}}\{\text{SNonce}, \text{sn}, \text{msg2}\}$

**[Message 3: A → S]**

AA, ANonce, sn+1, msg3,  $\text{MIC}_{\text{PTK}}\{\text{ANonce}, \text{sn}+1, \text{msg3}\}$

**[Message 4: S → A]**

SPA, sn+1, msg4,  $\text{MIC}_{\text{PTK}}\{\text{sn}+1, \text{msg4}\}$



```
C:\WINDOWS\system32\cmd.exe - aircrack.exe -w dic -0 wpa.cap
aircrack 2.3

[00:02:53] 24094 keys tested (138.77 k/s)

KEY FOUND! [ checkpassword ]

Master Key : B9 57 19 56 6F EB DF F6 0C 9E FD 84 F9 EB 10 F4
B8 8D AA 2F 41 FC D4 02 56 A1 2B CB B6 08 18 B5

Transient Key : 7F CD FA 92 14 B6 5C F1 F8 7D BC 9C 05 D8 CA 92
73 72 40 9E CD D7 CC 6D F5 A1 4D 58 10 15 04 B4
B6 5C 92 E5 AC CB 03 96 01 DD FA 4C C0 F4 8A 1F
C3 4F CA C5 3C 8A 09 8D 24 BB 42 0E C9 1F 97 9B

EAPOL HMAC : CB FB 97 81 DC 7E 41 B8 6B A1 48 B6 AC CE 4E 00

Press Ctrl-C to exit.
```

- **Password cracker (GPLv2 Licence)**
  - <http://www.openwall.com/john/>
  - Support for DES, MD5, Blowfish and others
- **Single Crack Mode**
  - Uses correlation with “username field”
- **Wordlist mode**
  - Dictionary password attack
    - Dictionaries at <ftp://ftp.ox.ac.uk/pub/wordlists>
- **Incremental Mode**
  - it can try all character combinations as passwords
- **External Mode**
  - heavily programmable

- **Syntax based on Crack 5.0a**
  - `ftp://ftp.cert.dfn.de/pub/tools/password/Crack/`
- **Examples:**
  - `-c[rules]`: ignore this rule if password are not case sensitive
  - `> 3`: do no process di word of smaller than 3 characters
  - `/`: convert alla characters to lowercase